

Proceedings



Md. Humayun Kabir\*, Md. Ahasan Kabir, Md. Saiful Islam, Mohammad Golam Mortuza, and Mohammad Mohiuddin

- <sup>1</sup> Department of Electronics and Telecommunication Engineering, Chittagong University of Engineering and Technology (CUET), Chittagong - 4349, Bangladesh; <u>mdhkrrabby@gmail.com</u> (M.H. K); <u>kabir.ece07@gmail.com</u> (M.A.K); <u>saiful05eee@cuet.ac.bd</u> (M.S.I); <u>mortuzaete@gmail.com</u> (M.G.M); <u>Muhin.cuet.bd@gmail.com</u> (M.M)
- \* Correspondence: mdhkrrabby@gmail.com; Tel.: +8801515286984
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Abstract: Computer network communication is quickly growing in this pandemic situation. Phone 13 Conferencing, Video Streaming, and Sharing File / Printing are all made easier with communica-14 tions technologies. Data transmitted in time with little interruption become a significant achieve-15 ment of wireless sensor networks (WSNs). A massive network is interconnection computer net-16 works in the globe connected by the internet and the internet plays a critical role in WSNs. Data 17 access is a key element of any enterprise network and the routing protocol is used to transmit data 18 or access data. Due to the growing use of WSNs, it is essential to know about the network structure, 19 the routing protocol. The routing protocols must be used to route all data sent over the internet 20 between the source and the destination. Which chooses the optimum routes between any two nodes 21 in an enterprise network. This research focused on how the routing table will determine the opti-22 mum path/route of data packets to be transmitted from source to destination. The performance of 23 three routing protocols, Routing Information Protocol (RIP), Enhanced Interior Gateway Routing 24 Protocol (EIGRP), and Open Shortest Path First (OSPF), is investigated in this research for the mas-25 sive mesh-based enterprise wireless sensor network. We also investigated the behaviors of end-to-26 end packet latency, convergence time on flapping connections, and average point-to-point through-27 put (bits/sec) between net-work links. Finally, the simulation results are compared to the efficacy 28 and performance of these protocols implemented in the Wireless LAN and internet based Wireless 29 Sensor Network. 30

**Keywords:** Wireless Sensor Network, RIP, EIGRP, OSPF, Redistribution, Dynamic Routing Protocol

## 1. Introduction

Today communication based on the internet has become an integral element of daily 34 life. Consequently, the development of computer networks based on IP routing proto-cols 35 also plays a significant role in any enterprise network. In small areas of a wireless sensor 36 network, the sensor node and base stations are so close to each other and share infor-37 mation directly with one another with low latency [1]. This type of communication system 38 is known as a single-hop communication system. But in wireless sensor networks, the 39 coverage area is massive, and to cover these large areas need many sensor nodes. Multi-40hop communication (indirect connection) is required in this case because most of the sen-41 sor nodes are locating so far away from the sensor nodes that they cannot connect directly 42 with the base station [2]. Within a multi-hop network connection, the sensor nodes not 43 only generate and transmit data but also provide a route for other sensor nodes to com-44 municate with the destination base station node. Finding an appropriate route from a 45 source node to a destination node is referred to as routing, and it is the key responsibility 46 of the routing protocol [3]. Routing protocols describe how routers interact with another 47 router, execute this task and identify the optimal paths for transferring data from one 48node to another. Data transmitted over the internet should be routed be-tween networks 49

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using routing algorithms [4–6]. The routing algorithms depend on different parameters 1 for selecting the most suitable path for transmitting information over the internet (i.e., 2 Bandwidth, cost, packet delay, hop count, and maximum transmission unit). The signifi-3 cant benefit of adopting a dynamic routing protocol allows routers to learn about new 4 networks when changes in the network topology and discover alternative routes if a link 5 fails in an existing network [7–10]. Currently, many organizations are shifting towards 6 their previous network topology (such as RIP) to more current net-work topology by up-7 grading routing protocols mechanisms (such as EIGRP, OSPF) [11–13]. This research aims 8 to identify the optimal routing protocol topology for each routing protocol since each rout-9 ing protocol has unique characteristics. Routing protocols define how routers acquire net-10 work topology information. Identifying the route should be accomplished using a pack-11 age router function that analyzes all possible routes to the destination and determines the 12 most optimum [14–16]. Each routing protocol facilitates the exchange of network infor-13 mation between participating routers. However, some proto-cols only convey information 14 about direct connections. There are also variations in speed and scalability across proto-15 cols [17–18]. 16

This research will compare the performance of several inner gateway dynamic rout-17 ing protocols, including RIPv2, EIGRP, and OSPF, using Cisco packet tracer soft-ware. In 18 addition, to demonstrate how to transfer data across various networks running different 19 routing protocols using route redistribution systems in packet tracer simulation software. 20 Each of these dynamic routing protocols has its own set of advantages and disadvantages; 21 for example, one protocol has rapid convergence, while another has high reliability. How-22 ever, the dynamic routing is all improved in general scalability, robust-ness, and conver-23 gence. 24

This work has three parts. First, a theoretical analysis of the three routing methods will be presented. Second, we'll look at how to create and execute a standard model for testing routing protocols. The simulation was done CISCO PACKET TRACER network simulation tool. Finally, we'll look at some of the outcomes and check our conclusions about them.

## 2. Concept of the Mesh-based Enterprise Wireless Sensor Network

It is a core network design with numerous redundant connections between network 31 nodes. Thus, if any nodes fail in a wireless mesh architecture, there are many alternative 32 ways to communicate with each node. A mesh network combines other topologies such 33 as Star, Ring, and Bus in a hybrid topology. Also, specific WAN architectures, like the 34 Internet, use mesh routing, which works even in disaster [1]. 35



Figure 1. Mesh-based Enterprise Wireless Sensor Network.

Full mesh and partial mesh are the two mesh topologies. When every node in a net-38work has a connection linking it to every other node, this is referred to as full mesh topol-39ogy. The maximum concentration of redundancy is achieved by using a full mesh. As a40result, if a node fails, network traffic may be diverted to any other node. Specific nodes41are fully meshed with partial mesh, while others are linked to one or two others. In peripheral networks connected to a full mesh backbone, the partial mesh is prevalent. Partial43mesh is less costly but less redundant than complete mesh [18].44

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## 3. Problem Description and Main Contribution

The primary objective of this research is to evaluate and compare the proposed rout-2 ing protocols using various performance metrics. This evaluation is carried out both the-3 oretically and through simulation. Routing is the process of transferring data from one 4 source to a destination. Typically, this data is routed through a series of intermediary de-5 vices. The objective of the routing protocols is to give the information necessary for send-6 ing the packet appropriately to these intermediate devices. Therefore, the routing proto-7 cols are essential in ensuring prevent network devices from connecting with one another. 8 Every routing protocol has an algorithm, and this algorithm must define techniques for 9 routing protocols to work correctly. Simulates networks using the CISCO Packet Tracer 10 8.0 simulator. 11

The conventional procedures are:

- $\checkmark$ These steps are used to receive and send network information.
- Second, finding the best route to a location and adding it to the routing data-base.

 $\checkmark$ Finally, it is a procedure to identifying, responding, and notifying network changes. As a result, different algorithms can affect total network performance. Thus, these significant research achievements are:

- $\checkmark$ To create two network topologies with RIP, EIGRP, and OSPF to analyze their performance.
- To simulate various network topologies using packets, transfer, and observe the performance differences between the OSPF, EIGRP, and RIP networks.
- Summarize and analyze the simulated findings.

## 4. Scenario Analysis

Figure 2 shows a full-mesh topology with OSPF, EIGRP, and RIP routing protocols. 24



Figure 2. Proposed Mesh-based Simulation Enterprise Wireless Sensor Network.

The next step is to test the system after all the settings and do the configurations. The 27 network topology is testing the scenario shown in figure 3. 28



Figure 3. Testing Process of Mesh Network Topology.

The underlying structure created a complete lattice geography reconstruction using 31 ten routers; each router directly connected to its neighbors' routers. Additionally, each of 32 these routers will serve a single client, as seen in Table 1. 33 34

| Router – 1 |             | Router – 2 |              | Router – 3 |              | ŀ    | Router – 4   | Router – 5 |              |  |
|------------|-------------|------------|--------------|------------|--------------|------|--------------|------------|--------------|--|
| G0/0       | 193.169.1.1 | G0/0       | 193.169.11.1 | G0/0       | 193.169.20.1 | G0/0 | 193.169.30.1 | G0/0       | 193.169.40.1 |  |
| G1/0       | 193.169.2.1 | G1/0       | 193.169.2.2  | G1/0       | 193.169.3.2  | G1/0 | 193.169.4.2  | G1/0       | 193.169.5.2  |  |
| G2/0       | 193.169.3.1 | G2/0       | 193.169.12.1 | G2/0       | 193.169.12.2 | G2/0 | 193.169.13.2 | G2/0       | 193.169.14.2 |  |
| G3/0       | 193.169.4.1 | G3/0       | 193.169.13.1 | G3/0       | 193.169.22.1 | G3/0 | 193.169.22.2 | G3/0       | 193.169.23.2 |  |

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| G4/0 | 193.169.5.1   | G4/0 | 193.169.14.1  | G4/0 | 193.169.23.1  | G4/0 | 193.169.31.1  | G4/0 | 193.169.31.2  |
|------|---------------|------|---------------|------|---------------|------|---------------|------|---------------|
| G5/0 | 193.169.6.1   | G5/0 | 193.169.15.1  | G5/0 | 193.169.24.1  | G5/0 | 193.169.32.1  | G5/0 | 193.169.41.1  |
| G6/0 | 193.169.7.1   | G6/0 | 193.169.16.1  | G6/0 | 193.169.25.1  | G6/0 | 193.169.33.1  | G6/0 | 193.169.42.1  |
| G7/0 | 193.169.8.1   | G7/0 | 193.169.17.1  | G7/0 | 193.169.26.1  | G7/0 | 193.169.34.1  | G7/0 | 193.169.43.1  |
| G8/0 | 193.169.9.1   | G8/0 | 193.169.18.1  | G8/0 | 193.169.27.1  | G8/0 | 193.169.35.1  | G8/0 | 193.169.44.1  |
| G9/0 | 193.169.10.1  | G9/0 | 193.169.19.1  | G9/0 | 193.169.28.1  | G9/0 | 193.169.36.1  | G9/0 | 193.169.45.1  |
| PC   | 193.169.1.10  | PC   | 193.169.11.10 | PC   | 193.169.20.10 | PC   | 193.169.30.10 | PC   | 193.169.40.10 |
| F    | Router – 6    | F    | Router – 7    | F    | Router – 8    | F    | Router – 8    | R    | outer – 10    |
| G0/0 | 193.169.50.1  | G0/0 | 193.169.60.1  | G0/0 | 193.169.70.1  | G0/0 | 193.169.80.1  | G0/0 | 193.169.90.1  |
| G1/0 | 193.169.6.2   | G1/0 | 193.169.7.2   | G1/0 | 193.169.8.2   | G1/0 | 193.169.9.2   | G1/0 | 193.169.10.2  |
| G2/0 | 193.169.15.2  | G2/0 | 193.169.16.2  | G2/0 | 193.169.17.2  | G2/0 | 193.169.18.2  | G2/0 | 193.169.19.2  |
| G3/0 | 193.169.24.2  | G3/0 | 193.169.25.2  | G3/0 | 193.169.26.2  | G3/0 | 193.169.27.2  | G3/0 | 193.169.28.2  |
| G4/0 | 193.169.32.2  | G4/0 | 193.169.33.2  | G4/0 | 193.169.34.2  | G4/0 | 193.169.35.2  | G4/0 | 193.169.36.2  |
| G5/0 | 193.169.41.2  | G5/0 | 193.169.42.2  | G5/0 | 193.169.43.2  | G5/0 | 193.169.44.2  | G5/0 | 193.169.45.2  |
| G6/0 | 193.169.51.1  | G6/0 | 193.169.51.1  | G6/0 | 193.169.52.2  | G6/0 | 193.169.53.2  | G6/0 | 193.169.54.2  |
| G7/0 | 193.169.52.1  | G7/0 | 193.169.61.1  | G7/0 | 193.169.61.2  | G7/0 | 193.169.61.2  | G7/0 | 193.169.63.2  |
| G8/0 | 193.169.53.1  | G8/0 | 193.169.62.1  | G8/0 | 193.169.71.1  | G8/0 | 193.169.71.2  | G8/0 | 193.169.72.2  |
| G9/0 | 193.169.54.1  | G9/0 | 193.169.63.1  | G9/0 | 193.169.72.1  | G9/0 | 193.169.81.1  | G9/0 | 193.169.81.2  |
| PC   | 193.169.50.10 | PC   | 193.169.60.10 | PC   | 193.169.70.10 | PC   | 193.169.80.10 | PC   | 193.169.90.10 |
|      |               |      |               |      |               |      |               |      |               |

According to table 1, a topology of RIP, EIGRP, and OSPF network and router configuration on route 1 to 10. Next the IP client on the other PC is chosen. The ping command used on each network to check for network connectivity in the following experiment, which is completed successfully. Afterward, the process of transmitting data packets from one network to another is carried out via traceroute, as seen in Figure 4a, 4b & 4c. 6

| (a)   | (b)  | (c)  |  |  |  |
|---|--|--|--|--|--|
| frace complete.   | Trace complete.  | Trace complete.  |  |  |  |
| 1 0 ms 0 ms 0 ms 193.169.1.1<br>2 0 ms 0 ms 0 ms 193.169.10.2<br>3 0 ms 0 ms 0 ms 193.169.90.10   | 1 0 ms 0 ms 0 ms 193.169.1.1<br>2 0 ms 0 ms 0 ms 193.169.10.2<br>3 0 ms 0 ms 0 ms 193.169.90.10  | 1 0 ms 0 ms 0 ms 193.169.1.1<br>2 0 ms 0 ms 0 ms 193.169.90.1  |  |  |  |
| fracing route to 193.169.90.10 over a maximum of 30 hops:   | Tracing route to 193.169.90.10 over a maximum of 30 hops:  | Tracing route to 193.169.90.1 over a maximum of 30 hops  |  |  |  |
| :\>tracert 193.169.90.10  | C:\>tracert 193.169.90.10  | C:\>tracert 193.169.90.1   |  |  |  |
| <pre>Ping statistics for 193.169.90.10:<br/>Packets: Sent = 4, Received = 4, Lost = 0 (0% loss),<br/>pproximate round trip times in milli-seconds:<br/>Minimum = Oms, Maximum = Oms, Average = Oms</pre>  | <pre>Ping statistics for 193.169.90.10:<br/>Packets: Sent = 4, Received = 4, Lost = 0 (0% Loss),<br/>Approximate round trip times in milli-seconds:<br/>Minimum = 0ms, Maximum = 0ms, Average = 0ms</pre>  | <pre>Ping statistics for 193.169.90.1:<br/>Packets: Sent = 4, Received = 4, Lost = 0 (0% loss),<br/>Approximate round trip times in milli-seconds:<br/>Minimum = Oms, Maximum = Oms, Average = Oms</pre>   |  |  |  |
| <pre>leply from 193.169.90.10: bytes=32 time<lms ttl="126&lt;br">leply from 193.165.90.10: bytes=32 time<lms ttl="126&lt;br">leply from 193.165.90.10: bytes=32 time<lms ttl="126&lt;br">leply from 193.169.90.10: bytes=32 time<lms ttl="126&lt;/pre"></lms></lms></lms></lms></pre> | Reply from 193.169.90.10: bytes=32 time <lms ttl="126&lt;br">Reply from 193.169.00.10: bytes=32 time<lms ttl="126&lt;br">Reply from 193.169.90.10: bytes=32 time<lms ttl="126&lt;br">Reply from 193.169.90.10: bytes=32 time<lms ttl="126&lt;/td"><td colspan="4">Reply from 193.169.90.1: bytes=32 time<lms ttl="254&lt;br">Reply from 193.169.90.1: bytes=32 time<lms ttl="254&lt;br">Reply from 193.169.90.1: bytes=32 time<lms ttl="254&lt;br">Reply from 193.169.90.1: bytes=32 time<lms ttl="254&lt;/td"></lms></lms></lms></lms></td></lms></lms></lms></lms> | Reply from 193.169.90.1: bytes=32 time <lms ttl="254&lt;br">Reply from 193.169.90.1: bytes=32 time<lms ttl="254&lt;br">Reply from 193.169.90.1: bytes=32 time<lms ttl="254&lt;br">Reply from 193.169.90.1: bytes=32 time<lms ttl="254&lt;/td"></lms></lms></lms></lms> |  |  |  |
| inging 193.169.90.10 with 32 bytes of data:   | Pinging 193.169.90.10 with 32 bytes of data:   | Pinging 193,169.90.1 with 32 bytes of data:  |  |  |  |
| :\>ping 193.169.90.10   | C:\>ping 193.169.90.10   | C:\>ping 193.169.90.1  |  |  |  |

**Figure 4a.** Packet Sending and Tracert Checking on Route 1 to Route 10 using RIP, **4b.** Packet Sending and Tracert Checking on Route 1 to Route 9 using EIGRP, **4c.** Packet Sending and Tracert Checking on Route 1 to Route 8 using OSPF.

# 5. Time Testing

Time testing is done after line termination when transmitting packets. The following table12summarizes the findings of the time tests performed in Table 2 for full-mesh routing using13RIP, EIGRP, and OSPF. First, conduct continuous ping tests, then pause; a delay time will14display. In addition, table 2 summarizes the experimental findings for routing RIP, EIGRP,15and OSPF on the mesh's entire topology.16

Table 2. Router Full Mesh RIP / EIGRP / OSPF.

| Full Mesh |           |      |       |      |           |           |      |       |      |  |  |
|-----------|-----------|------|-------|------|-----------|-----------|------|-------|------|--|--|
| Client PC | Client PC | RIP  | EIGRP | OSPF | Client PC | Client PC | RIP  | EIGRP | OSPF |  |  |
| PC – 1    | PC – 2    | 8 ms | 1 ms  | 2 ms | PC – 1    | PC – 7    | 8 ms | 9 ms  | 4 ms |  |  |

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| PC – 1 | PC – 3 | 6 ms | 7 ms | 2 ms | PC – 1 | PC – 8  | 9 ms  | 9 ms | 9 ms |
|--------|--------|------|------|------|--------|---------|-------|------|------|
| PC – 1 | PC – 4 | 8 ms | 8 ms | 3 ms | PC – 1 | PC – 9  | 11 ms | 9 ms | 5 ms |
| PC – 1 | PC – 5 | 8 ms | 8 ms | 5 ms | PC – 1 | PC – 10 | 12 ms | 8 ms | 6 ms |
| PC – 1 | PC – 6 | 7 ms | 8 ms | 6 ms | -      | -       | -     | -    | -    |

Table 3 below shows half-mesh time results for RIP, EIGRP, and OSPF routing. **Table 3.** Router Half Mesh RIP / EIGRP / OSPF.

| Half Mesh |           |       |       |      |           |           |       |         |      |  |
|-----------|-----------|-------|-------|------|-----------|-----------|-------|---------|------|--|
| Client PC | Client PC | RIP   | EIGRP | OSPF | Client PC | Client PC | RIP   | EIGRP   | OSPF |  |
| PC – 1    | PC – 2    | 12 ms | 3 ms  | 2 ms | PC – 1    | PC – 7    | 8 ms  | 9.66 ms | 5 ms |  |
| PC – 1    | PC – 3    | 7 ms  | 2 ms  | 2 ms | PC – 1    | PC – 8    | 9 ms  | 9.33 ms | 3 ms |  |
| PC – 1    | PC – 4    | 8 ms  | 5 ms  | 3 ms | PC – 1    | PC – 9    | 9 ms  | 9.66 ms | 3 ms |  |
| PC – 1    | PC – 5    | 9 ms  | 4 ms  | 3 ms | PC – 1    | PC – 10   | 10 ms | 8.33 ms | 5 ms |  |
| PC – 1    | PC – 6    | 8 ms  | 8 ms  | 4 ms | -         | -         | -     | -       | -    |  |

#### 6. Analysis Results

A simulation duration of four minutes for voice, HTTP, and video data transport is specified for LAN-to-server and server-to-LAN configurations in full and half mesh RIP, 5 OSPF, and EIGRP. Figure 5a shows the average voice packet end-to-end latency topology. 6



**Figure 5a.** The Average Voice Packet end-to-end Latency, **5b.** The Average Value of HTTP Page Response Time, **5c.** The Average Video Packet end-to-end Latency.

Figure 5b depicts the response time of an HTTP page for a simulated network. Based on distance-vector techniques, the RIP routing protocol showed better performance than other routing protocols. OSPF performs better in video transfer, responds faster to network changes, and better utilizes bandwidth, resulting in a minimal delay, as seen in figure 5c. 14



Figure 6. The Average Point to Point Throughput (bit/sec).

Figure 6 illustrates the average network throughput for three protocols. The OSPF protocols produce better throughput than any of the other protocols evaluated in this test case. The following result is the findings of the analysis based on the experiments.

#### 7. Conclusion

In a wireless communication network, identify the optimal path from the sensor node 21 to the destination is more difficult. The routing protocols help to find an optimal path be-22 tween source and destination nodes and minimize these difficulties. The optimal path se-23 lection depends upon several factors. This research discusses and analyzes different para-24 metric aspects of routing protocols. RIP, EIGRP, and OSPF routing protocols were analyzed 25 and evaluated via an extensive simulation process using carefully selected parameters to 26 acquire the features of their routing algorithms. The measured metrics are voice, HTTP, and 27 video traffic transmitted and received, as well as average end-to-end latency and average 28 point-to-point throughput. The protocol RIP has shown the most significant uncertainty, 29

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whereas OSPF has demonstrated the lowest latency. Furthermore, the OSPF protocols attain a better throughput than any other protocols evaluated in this test scenario. From the above result, we see that the OSPF routing protocol is more suitable for multi-hop wireless sensor networks.

In future research, we will work on simulations with much more realistic topologies and increased optimization accuracy to enhance and show the efficacy of routing protocols in terms of wireless sensor network performance.

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