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## Efficient numerical evaluation of weak restricted compositions

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<p><b>Graphical Abstract</b></p> <div style="text-align: center;"> </div> <p>Combinations with Repetition <math>C_{(n,m)} = \frac{(n+m-1)!}{m!(n-1)!}</math></p> <table style="width: 100%; text-align: center;"> <tr> <td>1. </td> <td>8. </td> <td>15. </td> </tr> <tr> <td>2. </td> <td>9. </td> <td>16. </td> </tr> <tr> <td>3. </td> <td>10. </td> <td>17. </td> </tr> <tr> <td>4. </td> <td>11. </td> <td>18. </td> </tr> <tr> <td>5. </td> <td>12. </td> <td>19. </td> </tr> <tr> <td>6. </td> <td>13. </td> <td>20. </td> </tr> <tr> <td>7. </td> <td>14. </td> <td></td> </tr> </table>	1.	8.	15.	2.	9.	16.	3.	10.	17.	4.	11.	18.	5.	12.	19.	6.	13.	20.	7.	14.		<p><b>Abstract.</b></p> <p><i>We propose an algorithm to calculate the number of weak compositions, wherein each part is restricted to a different range of integers. This algorithm performs different orders of approximation up to the exact solution by using the Inclusion-Exclusion Principle. The great advantage of it with respect to the classical generating function technique is that the calculation is exponentially faster as the size of the numbers involved increase</i></p>
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### Introduction

If we want to count the different ways of gathering  $n$  elements from  $m$  different types in such a way we have  $n_k$  elements of type  $k = 1, \dots, m$ , i.e.

$$n = n_1 + n_2 + \dots + n_m, \tag{1}$$

we use the formula of combinations with repetition [3, Eqn. 5.2]

$$M = \binom{n+m-1}{m-1}. \tag{2}$$

However, (2) is based on the fact that there are available as many elements of any type  $k$  as we need. Therefore, a natural generalization of this problem is to consider that each  $n_k$  is bounded by lower and upper limits. According to this generalization, define  $n_{\min,k}$  and  $n_{\max,k}$  as the minimum and maximum number of items we can select of type  $k$ , i.e.  $n_{\min,k} \leq n_k \leq n_{\max,k}$ . Also, define  $M(n, \vec{n}_{\min}, \vec{n}_{\max})$  as the number of integral solutions of (1), where  $\vec{n}_{\min} = (n_{\min,1}, \dots, n_{\min,m})$  and  $\vec{n}_{\max} = (n_{\max,1}, \dots, n_{\max,m})$ .

To compute  $M(n, \vec{n}_{\min}, \vec{n}_{\max})$ , we can consider the approach given in [1, Sect. 6.2], which uses the combination of repetition formula (2) as well as the Inclusion-Exclusion Principle [2, p.177]. In [4], this

approach is modified in order to obtain an efficient algorithm to compute  $M(n, \vec{n}_{\min}, \vec{n}_{\max})$ . Next, we present the MATHEMATICA program to compute  $M(n, \vec{n}_{\min}, \vec{n}_{\max})$  with the latter approach.

## Materials and Methods

```
(* Auxiliary functions *)
fk[ns_,sv_,m_]:=Module[{tk,k},
tk=Total[sv];
k=Length[sv];
If[ns>tk+k-1,Binomial[ns-tk+m-k-1,m-1],0]
];
Sk[ns_,m_,v_,k_]:=Module[{ind,l,s,i},
l=Length[v];
ind=Subsets[Range[l],{k}];
s=0;
For[i=1,i<=Length[ind],i++,
s+=fk[ns,v[[ind[[i]]]],m];
];
s
];
(* Main function (from group k, we can select 0 to v[[k]] items *)
M0[n_,v_]:=Module[{w,m,t,ns,h},
w=Sort[v];
m=Length[v] (* Number of groups *);
t=Total[v] (* Number of elements we can choose *);
ns=Min[n,t-n] (* We set the direct (take elements) or the complementary (neglect elements) problem *);
(* Calculation of the order *)
h=0;
While[Total[Take[w,h+1]]+h+1-ns<=0,h++];
{h,Sum[(-1)^k Sk[ns,m,v,k],{k,0,h}]}
];
(* Generalization in which we have to select the number of items of each group between a minimum and a maximum *)
M1[n_,vmin_,vmax_]:=M0[n-Total[vmin],vmax-vmin];
```

## Results and Discussion

The great advantage of the code based on the Inclusion-Exclusion Principle is that it is much faster than the one based on the generating function. This is so because the Inclusion-Exclusion Principle code does not involve any symbolic processing. Fig. 1 shows the time ratio performance between both codes as a function of the size of the numbers involved in the calculations. This size is a scale factor, i.e. size 8 means that on average the numbers are double than size 4. Size 1 means that we consider 1 digit numbers. It is apparent that the time ratio performance increases exponentially with size, being the Inclusion-Exclusion Principle code much more efficient. To compute the calculations of Fig. 1, we have set  $m = 8$ , but similar patterns can be found for other values of  $m$ .

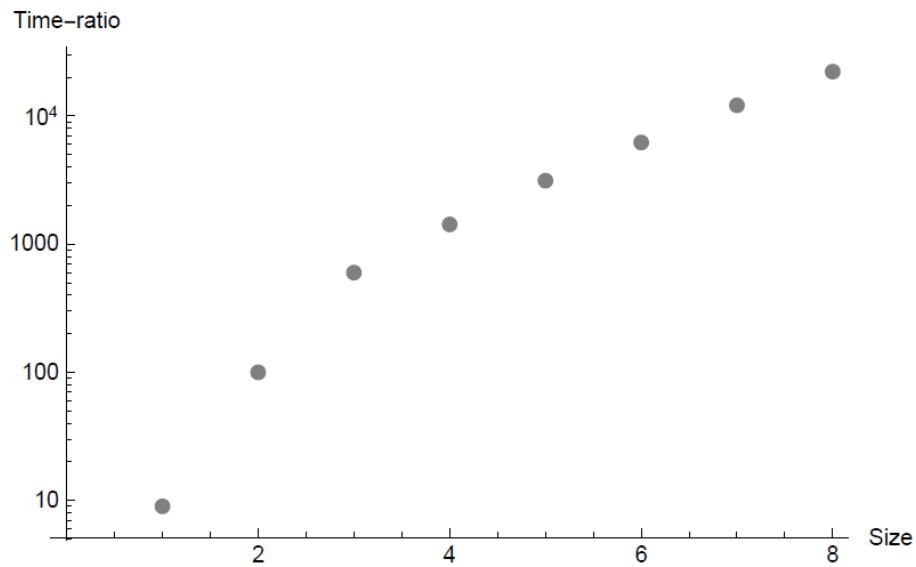


Figure 1. Time ratio between the Inclusion-Exclusion Principle coder and the generating code.

### References

- [1] Brualdi RA. Introductory Combinatorics. Pearson Education India; 1977.
- [2] Comtet L. Advanced Combinatorics: The art of infinite expansions. Springer Science & Business Media, 2012.
- [3] Feller W. An introduction to probability theory and its applications, vol. 2. John Wiley & Sons, 2008.
- [4] González-Santander, Juan Luis. "Efficient numerical evaluation of weak restricted compositions." *Nereis. Interdisciplinary Ibero-American Journal of Methods, Modelling and Simulation*. 14 (2022). DOI: 10.46583/nereis\_2022.14.1018