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Optimal Scheduling for Precedence-Constrained Applications on Heterogeneous Machines

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Outline

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Introduction: High performance computing

- High Performance Computing(HPC) systems are fundamental in a wide variety of scenarios as Wall Street or NASA [1], HPC systems requires high amount of power [2],
- HCP systems also known as super computers are formed by a set of machines interconnected over a bus or network platform.



Figure 1: Photo taken on June 16, 2013 shows the supercomputer Tianhe-2 developed by China's National University of Defense Technology. ([tv.81.cn/He Shuyuan](http://tv.81.cn/HeShuyuan))

Introduction: Problem formulation

Task	m_1	m_2	m_3	\bar{p}_i
0	11	13	9	11.0
1	10	15	11	12.0
2	09	12	14	11.6
3	12	16	10	12.3
4	15	11	19	15.0
5	13	09	05	8.60
6	11	15	13	12.3
7	11	15	10	12.0

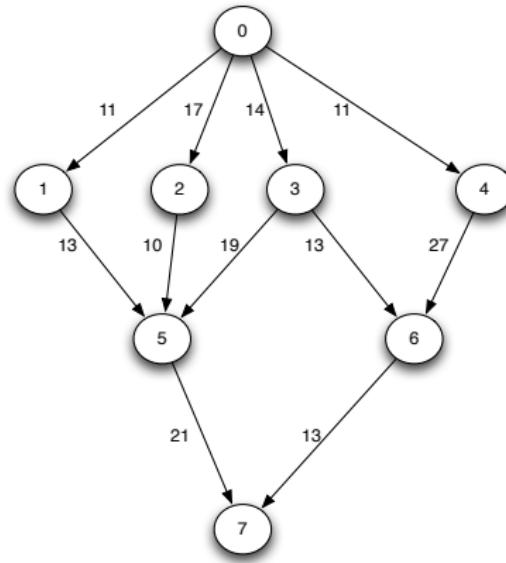


Figure 2: Example of a scheduling problem with a DAG.

Introduction: Problem definition

Given a DAG, $G = (T, E)$, where:

- T is a set of n tasks $T = \{t_1, t_2, \dots, t_n\}$.
- E is a set of edges, $(t_i, t_j) \in E$ are the precedence constraints.

A HCS is represented by:

- M is a set of machines $M = \{m_1, m_2, \dots, m_k\}$.
- P a processing time for each task i in each machine k .

$$ts_i = \max\{\max\{tf_j\} \forall (t_j, t_i) \in E, Pw_{i,k} \in m_k\} \quad (1)$$

$$tf_i = ts_i + P_{i,k} + C_{i,k} \quad (2)$$

$$Makespan \quad \max_{i \in \{t_1, \dots, t_n\}} tf_i \quad (3)$$

Methodology: Branch and bound

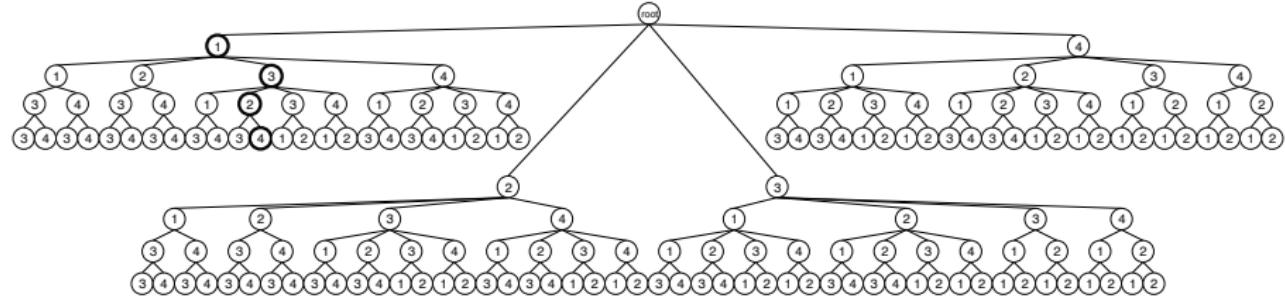


Figure 3: Example of a branch and bound tree.

Methodology: Branch and bound

Algorithm 1 Branch and bound scheduling.

```
1: function Scheduling( $S_{HEFT}$ )
2:    $S_{best} \leftarrow S_{HEFT}$ 
3:    $Bound \leftarrow \text{evaluate } S_{best}$ 
4:   Initialize queue  $Q$  pushing  $\{0, 1, \dots, n - 1\}$ 
5:    $TopologicalOrder(0)$ 
6:   return  $S_{best}$ 
```

Methodology: Branch and bound

Algorithm 2 Computation of the topological order.

```
1: function TopologicalOrder(deep)
2:   for i in deep, ..., n do
3:     e  $\leftarrow Q.pop()$ 
4:     ord[level]  $\leftarrow e$ 
5:     Evaluate partial topological order  $\{ord[0], ord[1], \dots, ord[deep + 1]\}$ 
6:     if partial order ord is feasible then
7:       if level < n - 1 then
8:         TopologicalOrder(deep + 1)
9:       else  $\triangleright$  The order is complete.
10:      MachineAssignment(0)
11:      Q.push(e)
```

Methodology: Branch and bound

Algorithm 3 Machine assignment.

```
1: function MachineAssignment(deep)
2:   for j in  $0, \dots, k$  do
3:     asig[ord[deep]]  $\leftarrow j$ 
4:     Scurrent  $\leftarrow (\textit{ord}, \textit{asig})$ 
5:     Evaluate partial solution Scurrent in range  $\{0, 1, \dots, x + 1\}$ 
6:     if partial evaluation < Bound then
7:       if prof < n – 1 then
8:         MachineAssignment(deep + 1)
9:       else  $\triangleright$  The solution is complete.
10:      Sbest  $\leftarrow \textit{S}_{\textit{current}}$ 
11:      Bound  $\leftarrow$  partial evaluation
```

Methodology: Mixed integer linear programming model

$$\min \quad z \quad (4)$$

Subject to:

$$\sum_{k=1}^m x_{ik} = 1, \quad i = 1, \dots, n. \quad (5)$$

$$f_{jk} \leq s_{ik} + M(3 - x_{ik} - x_{jk} - q_{ij}), \quad \forall j \prec i; \\ i, j = 1, \dots, n; k = 1, \dots, m; i \neq j \quad (6)$$

$$f_{jk_2} \leq s_{ik_2} + M(3 - x_{ik_2} - x_{jk_1} - q_{ij}) + c_{ij}, \quad \forall j \prec i; \\ i, j = 1, \dots, n; k_1, k_2 = 1, \dots, m; i \neq j; k_1 \neq k_2 \quad (7)$$

$$s_{ik} + P_{ik} \leq f_{ik} + (M - x_{ik}), \quad i = 1, \dots, n; k = 1, \dots, m \quad (8)$$

$$f_{jk} - s_{ik} \leq M(1 - x_{jk}), \quad \forall j \prec i; i, j = 1, \dots, n; k = 1, \dots, m \quad (9)$$

Methodology: Mixed integer linear programming model

$$f_{jk_2} - s_{ik_1} \leq M(1 - x_{jk_2}) + c_{ij}, \quad \forall j \prec i; i, j = 1, \dots, n; k = 1, \dots, m \quad (10)$$

$$f_{ik} \leq s_{jk} + M * y, \quad i, j = 1, \dots, n; k = 1, \dots, m; i \neq j \quad (11)$$

$$f_{jk} \leq s_{ik} + M(1 - y), \quad i, j = 1, \dots, n; k = 1, \dots, m; i \neq j \quad (12)$$

$$z \geq f_{ik}, \quad i = 1, \dots, n; k = 1, \dots, m \quad (13)$$

$$f_{ik} \geq s_{ik}, \quad i = 1, \dots, n; k = 1, \dots, m \quad (14)$$

$$0 \leq x_{ik} \leq 1, \quad 0 \leq y \leq 1, \quad 0 \leq f_{ik}, \quad 0 \leq s_{ik} \quad (15)$$

$M = \text{A big number}$

(16)

Experimentation setup

- A time limit of 37,297 seconds.
- A small instance set from diverse authors (See Table 1).

Results

Table 1: Time results for synthetic applications in the literature

Instance	Optimal	Reference	BB	Time	MILP	Time
Ahmad_3_9_28	22	[3]	22	0.13	22	0.46
Bittencourt_3_9_184	184	[4]	184	0.41	184	0.73
Hsu_3_10_84	80	[5]	80	14.36	80	0.60
Demiroz_3_7_47	47	[6]	47	0.00	47	0.13
Eswari_2_11_—	100	[7]	100	5.61	100	0.71
Eswari_2_11_61	56	[7]	56	3.69	56	1.52
Hamid_3_10_—	100	[8]	100	5.59	100	0.70
Ilavarasan_3_10_77	73	[9]	73	6.84	73	0.64
Ilavarasan_3_15_114	121	[10]	121	37297.23	121	28.05
Kang_3_10_76	73	[11]	73	7.88	73	2.34
Kang_3_10_84	79	[12]	79	17.80	79	0.76
Kuan_3_10_28	25	[13]	26	1.89	26	0.75
Liang_3_10_80	73	[14]	73	6.83	73	0.65
Linshan_4_9_38	40	[15]	40	11.31	40	1.99
Mohammad_2_11_64	56	[16]	56	3.83	56	1.51
SahA_3_11_131	118	[17]	118	6.88	118	3.43
SahB_3_6_76	66	[17]	66	0.00	66	0.13
Samantha_5_11_31	32	[18]	32	3696.81	32	6.21
Sample_3_8_100	81	[19]	81	0.08	81	0.75
Liang_3_10_80	73	[14]	73	6.83	73	0.67
YCLee_3_8_80	66	[20]	66	0.11	66	0.28
Average time in seconds				1956.86	2.52	

Conclusions

- Both strategies achieve the optimal values.
- The MILP gets a better computational time.
- A small instance set is proposed.

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